

# Learnable and Instance-Robust Predictions for Online Matching, Flows, and Load Balancing

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# ML is Everywhere

- Massive successes in image classification, NLP, etc.
- Widening array of applications
- Can ML help to improve algorithms for combinatorial problems?
  - (Online algorithms, in particular)

# Current Status

- Ski Rental
  - [Purohit et al. 2018], [Gollapudi et al. 2019]
- Caching
  - [Lykouris et al. 2018], [Rohatgi 2020]
- Scheduling
  - [Purohit et al. 2018], [Lattanzi et al. 2020]
- Secretaries
  - [Antoniadis et al 2020], [Dütting et al. 2021]
- ... and more



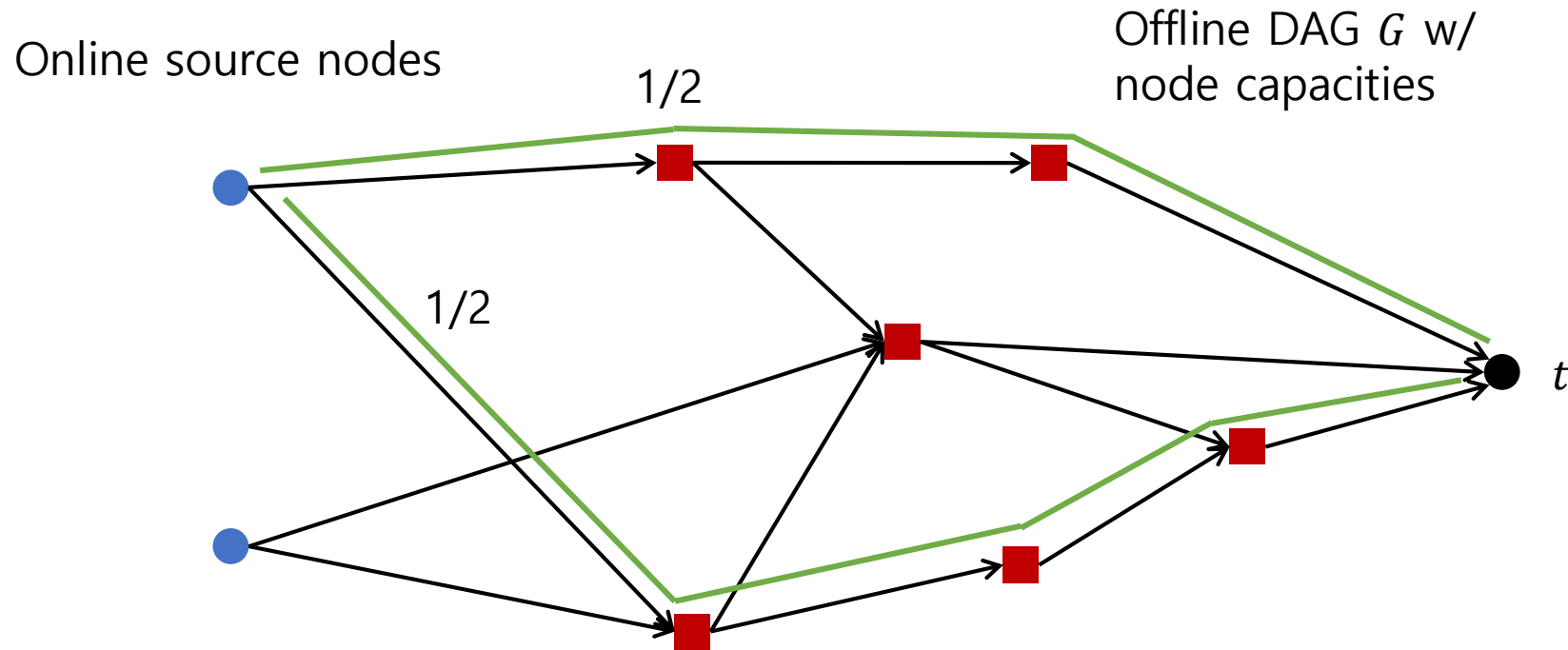
# Learning Augmented Algorithms

- Algorithm given access to predictions about online input
- There may be errors – algorithm must account for this
  - “Consistency” – no error case
  - “Robustness” – infinite error case
- Benefits
  - Accurate predictions can go beyond worst case
  - Focuses on algorithm design – abstracts the learning step

# Contributions

- Develop a new model for online algorithms with predictions
- Instance robustness
  - New way to understand “error”
- Learnability
  - What types of predictions are reasonable to construct?
- Focus on two online problems
  - Online Flow Allocation in DAGs
  - Online Load Balancing with Restricted Assignments

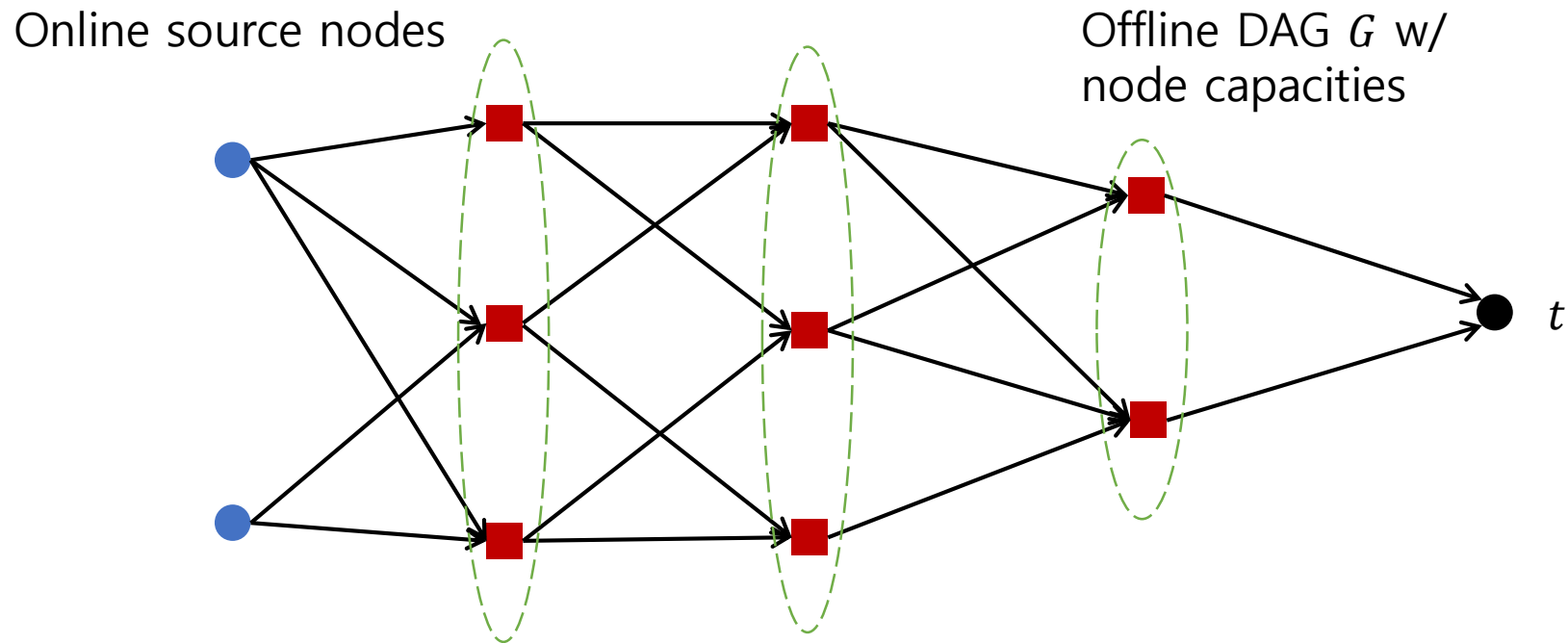
# Online Flow Allocation in DAGs



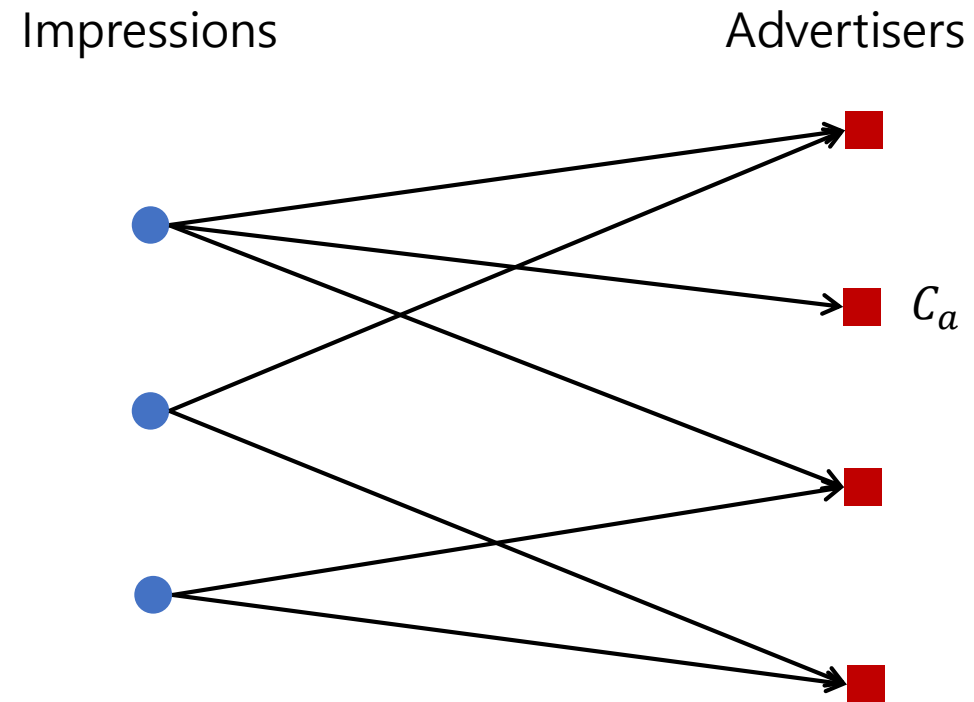
- Assign fractional unit flow to  $t$  from each online node  $s$  subject to node capacities
- Goal: maximize total value reaching  $t$

# Online Flow Allocation in DAGs

## Layered case

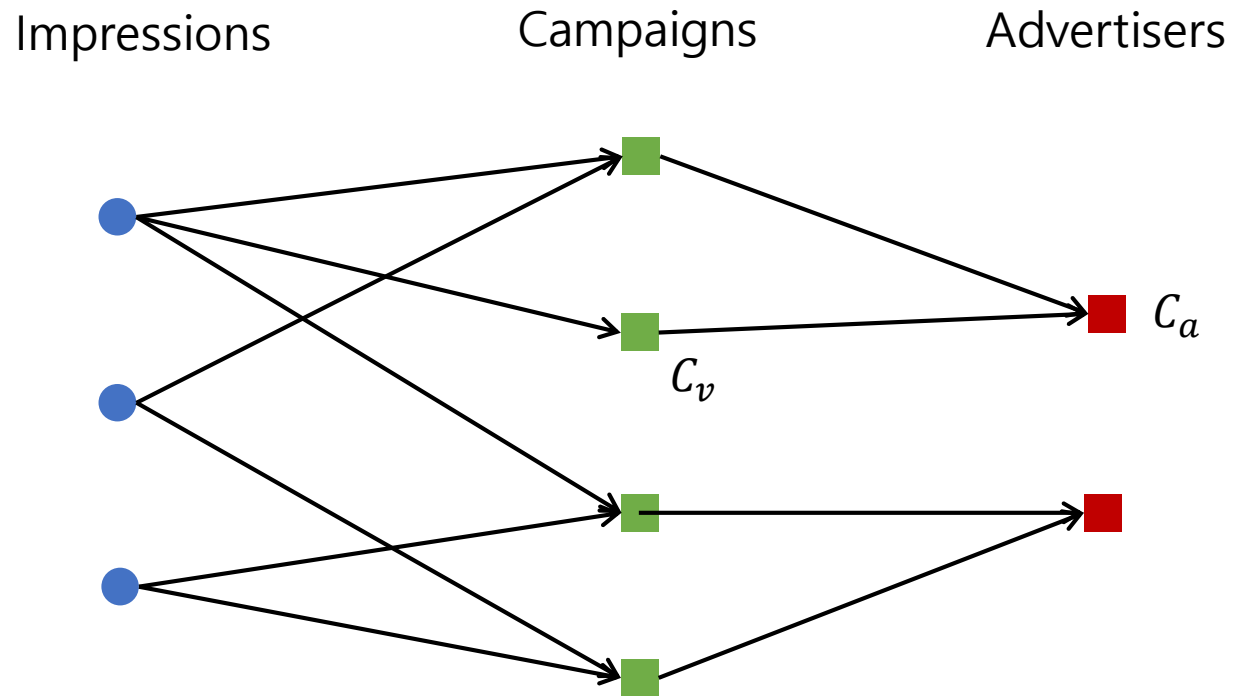


# Capacitated Online Matching





# Online Matching w/ Campaigns



# Parameter Robustness

- Algorithm given parameters  $\hat{y}$ 
  - Prediction of true parameters  $y^*$
  - E.g. guess of length of ski trip vs. true length
- Parameterize competitiveness by error  $\gamma$ 
  - E.g.  $\gamma = \|\hat{y} - y^*\|_1$
- Error typically depends on problem/prediction considered
- How to compare algorithms for the same problem with different types of predictions?

# Instance Robustness

- Instance  $\mathcal{I}$  is a vector of impression types
  - Type defined by outgoing neighborhood of an impression
  - $\mathcal{I}_i$  = number of impressions of type  $i$
- $y(\mathcal{I})$  = "correct" prediction for instance  $\mathcal{I}$
- Algorithm given  $y(\mathcal{I})$  as advice
- Sees instance  $\mathcal{I}'$  as online input
- Parameterize competitiveness by  $\gamma = \|\mathcal{I} - \mathcal{I}'\|_1$

# Learnability

- Formulate in terms of Data-driven algorithm design
  - [Gupta and Roughgarden, 2017], [Balcan et al. 2019]
  - Similar to PAC learning
- $ALG(I, y)$  = algorithm's value when given prediction  $y$
- Unknown distribution  $\mathcal{D}$  over instances
- Best prediction  $y^* = \arg \max_y \mathbb{E}_{I \sim \mathcal{D}}[ALG(I, y)]$
- Use samples from  $\mathcal{D}$  to compute  $\hat{y}$  s.t. following occurs w.h.p.:
$$\mathbb{E}_{I \sim \mathcal{D}}[ALG(I, \hat{y})] \geq (1 - \epsilon) \mathbb{E}_{I \sim \mathcal{D}}[ALG(I, y^*)]$$
  - Ideally polynomial number of samplers
  - Can also compare to  $\mathbb{E}_{I \sim \mathcal{D}}[OPT(I)]$

# Node Weights

- Predict a weight  $\alpha$  for each offline node
- Lower weight means more restrictive node
- Use weights to define a fractional flow
  - Node  $u$  has a weight  $\alpha_u$
  - For each edge  $(u, v)$  let  $x_{uv}(\alpha) = \frac{\alpha_v}{\sum_{v' \in N(u)} \alpha_{v'}}$
  - Node  $u$  sends an  $x_{uv}$ -fraction of its incoming flow along edge  $(u, v)$

# Existence of Good Weights

- Good weights give a  $(1 - \epsilon)$ -approximate flow
- Existence for matching case shown by [Agrawal, Zadimoghaddam, Mirrokni 2018]
- Extension to general DAGs requires significant work
  - Careful algorithm and analysis

# Results – Instance Robustness

- Instance  $\mathcal{I}$  is a vector of impression types
- Algorithm given good weights  $\alpha(\mathcal{I})$  as advice
- Sees instance  $\mathcal{I}'$  as online input
- Parameterize competitiveness by  $\gamma = \|\mathcal{I} - \mathcal{I}'\|_1$

Theorem: There is an online algorithm which achieves value

$$\max\left\{(1 - \epsilon)OPT - 2\gamma, \frac{1}{d + 1}OPT\right\}$$

where  $d$  is the diameter of  $G$  without  $t$

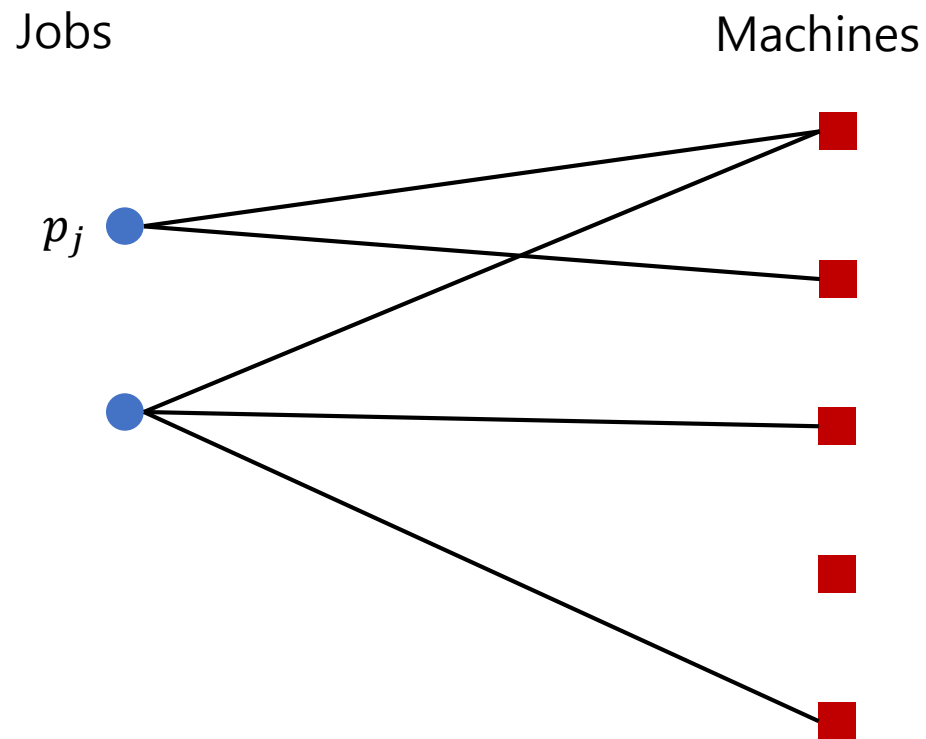
# Results - Learnability

- Show learnability under two assumptions
  - $\mathcal{D}$  is a product distribution (impressions are independent)
  - The optimal flow in the "expected instance" routes at least a constant amount of flow through each node

Theorem: Under the two assumptions, there is an algorithm with polynomial sample complexity for learning the weights



# Results – Load Balancing



- Consider node weights
- Existence of weights and parameter robustness considered in [Lattanzi et al. 2020]
- Show instance robustness and learnability under similar assumptions

# Conclusions and Future Work

- Consider a new model for online algorithms with predictions
  - Learnability + Instance robustness
- How to construct predictions from past data?
- How do we measure robustness?
- Demonstrate improvements empirically

Thank you! Questions?